```
choice-idempotency : ∀ {i : Size} {A : Domain} {D : Dimension} {e : BCC i A}
```

```
choice-idempotency {i} {A} {D} {e} = extensionality (\lambda \in A
            [ D ( e , e ) ] c
           <u>[ if (c D) then e</u>else e ] c \equiv (Eq.cong (flip [_] c) (if-idemp (c D))
                                                                                                                                                                                    •)
           [e]c
0 - 8.9k BCC.lagda.md Agda
completeness-by-expressiveness : \forall \{L_1, L_2 : VarLang\} \{C_1, C_2 : ConfLang\} \{S_1 : Set A = Set A =
           \rightarrow Complete L<sub>1</sub> C<sub>1</sub> S<sub>1</sub>
           \rightarrow L<sub>2</sub>, S<sub>2</sub> is-as-expressive-as L<sub>1</sub>, S<sub>1</sub>
           \rightarrow Complete L<sub>2</sub> C<sub>2</sub> S<sub>2</sub>
completeness-by-expressiveness {L<sub>1</sub>} {L<sub>2</sub>} {_} {S<sub>1</sub>} {S<sub>2</sub>} encode-in-L<sub>1</sub> L<sub>1</sub>-to-L<sub>2</sub>
          let {
🧿 – 11k Completeness.lagda.md Agda
```

#### Formal Languages for Solution-Space Variability

Paul Bittner, Jeffrey Young, Parisa Ataei, Alexander Schultheiß, Eric Walkingshaw, Leopoldo Teixeira, Thomas Thüm | FOSD 2023



Software Engineering Programming Languages































#### Using Formal Languages for Variability

Core Choice Calculus

Binary Choice Calculus

**Option Calculus** 

Artifact Trees

Algebraic Decision Diagrams

Variation Trees

Variability-Aware Abstract Syntax Trees



Algebraic Decision Diagrams

Variation Trees

Variability-Aware Abstract Syntax Trees



## Using Formal Languages for Variability





Binary Decision Diagrams

## Using Formal Languages for Variability





Binary Decision Diagrams





#### **Research Goal** Map Out Language Space







#### Why relating?

transfer research results





#### Bittner et al.







$$e ::= a \prec e, \dots, e \succ$$
 Object Structure  
 $| D\langle e, \dots, e \rangle$  Choice



$$e ::= a \prec e, \dots, e \succ$$
 Object Structure  
 $\mid D\langle e, \dots, e \rangle$  Choice





$$e ::= a \prec e, \dots, e \succ$$
 Object Structure  
 $\mid D\langle e, \dots, e \rangle$  Choice







$$e ::= a \prec e, \dots, e \succ$$
 Object Structure  
 $\mid D\langle e, \dots, e \rangle$  Choice





 $Salad? \langle \mathbf{v}, \circ \rangle,$   $Salad? \langle \mathbf{v}, \circ \rangle,$   $Patty \langle \mathbf{c}, \mathbf{o} \rangle, \mathbf{o} \rangle,$   $Sauce \langle \circ, \mathbf{c}, \mathbf{b}, \mathbf{c} \rangle,$   $\succ ]_{c}$ 



if 
$$c(Salad?) = 0$$
,  
 $c(Patty) = 0$ ,  
 $c(Sauce) = 2$ .













 $L_1$  is as expressive as  $L_2$ 

**iff** Every expression in  $L_2$  can be translated to an expression in  $L_1$  that describes the same set of variants.

 $L_1$  is as expressive as  $L_2$ 

**iff** Every expression in  $L_2$  can be translated to an expression in  $L_1$  that describes the same set of variants.

 $L_1$  is variant equivalent to  $L_2$ 

iff  $\begin{array}{c} L_1 \text{ is as expressive as } L_2 \\ \text{and vice versa.} \end{array}$ 

 $L_1$  is as expressive as  $L_2$ 

 $L_1$  is variant equivalent to  $L_2$ 

 $L_1$  is semantically equivalent to  $L_2$  iff

- Every expression in  $L_2$  can be translated to an expression in  $L_1$  that describes the same set of variants.
- **iff**  $L_1$  is as expressive as  $L_2$  and vice versa.
  - $L_1$  and  $L_2$  are variant equivalent and same configurations yield same variants. (Translation of configurations is an isomorphism.)



ł w	<b>Annotation Language</b> How to annotate elements <i>i</i> th variability information?	
	– higher-order logic	
	– propositional logic	
	+ list of literals	
	+ literals	
	– names	Composition
		How to derive variants?



	Annotatior How to anno with variability	<b>Language</b> tate elements y information?	
<u> </u>		- higher-order logic	
$$ Salad? $\langle \widehat{\mathbf{W}}, \circ  angle,$		propositional logic	
		- list of literals	
$Patty \langle \bigcirc , \bigcirc \rangle, \\ Sauce \langle \circ, \bigcirc, \land, \land, \land \rangle \rangle$		- literals	
>	Core Choice Calculus [Walkingshaw, 2013]	– names	Composition
-			<ul> <li>How to derive</li> </ul>
	Alternatives choose exactly one from a range of alternatives	<b>Options</b> in- or exclude an element	variants?

	Annotation How to anno with variability	<b>Language</b> tate elements y information?	
<b>—</b> ≺	-	- higher-order logic	
$\overbrace{Salad}^{Salad}$	-	- propositional logic	
🔗, Patty/ 🥽 🔊	-	- list of literals	
$Sauce(\circ, \bigcap, \blacktriangle, \bigcap, \Diamond)$	-	- literals	
>	Core Choice Calculus [Walkingshaw, 2013]	- Option Calculus	Composition
			- How to dorive
	Alternatives choose exactly one from a range of alternatives	<b>Options</b> in- or exclude an element	variants?





✓
 Salad?{♥},
 Ø,
 Tofu?{♥},
 Meat?{♥},
 Ketchup?{●},
 Mayo?{●}
 ✓

**Option Calculus** 



Named options cannot express alternatives!





Named options cannot express alternatives!





Variation Trees [Bittner et al., 2022] *L* is complete

iff *L* can encode any given set of variants.

# *L* is complete **iff** *L* can encode any given set of variants.

 $L_1$  is complete  $\land L_2$  is as expressive as  $L_1$ 

 $L_2$  is complete.

L is complete	iff	L can encode any given set of variants.
$L_1$ is complete $\land L_2$ is as expressive as $L_1$	Þ	$L_2$ is complete.
$L_1$ is complete $\land L_2$ is incomplete	Þ	$L_2$ is less expressive than $L_1$ .







- complete
- more expressive

KEN

229600 HI

 encoding options requires neutral domain elements or sacrifices sharing



229600 PUSH START

- complete
- more expressive

KEN

229600 HI

 encoding options requires neutral domain elements or sacrifices sharing





• incomplete

229600 PUSH START

less expressive

RYU

 cannot encode alternatives

- Options are useful syntax to increase sharing.
- For completeness, **else** statements or **negations** of annotations are essential.

Core Choice Calculus

Binary Choice Calculus

**Option Calculus** 

Artifact Trees

Algebraic Decision Diagrams

Variation Trees

Variability-Aware Abstract Syntax Trees

Binary Decision Diagrams

Core Choice Calculus		
Binary Choice Calculus	Option Calculus	Artifact Trees
Algebraic Decision Diagrams	Variation Trees	Variability-Aware Abstract Syntax Trees
Binary Decision Diagrams		

option calculus to formally clarify relationship between alternatives and options



option calculus to formally clarify relationship between alternatives and options

formal framework based on meta-language for variability



option calculus to formally clarify relationship between alternatives and options

formal framework based on meta-language for variability

formal comparison of variability languages



option calculus to formally clarify relationship between alternatives and options

formal framework based on meta-language for variability

formal comparison of variability languages

(in)completeness proofs



option calculus to formally clarify relationship between alternatives and options

formal framework based on meta-language for variability

formal comparison of variability languages

(in)completeness proofs

open-source Agda library



option calculus to formally clarify relationship between alternatives and options

formal framework based on meta-language for variability

formal comparison of variability languages

(in)completeness proofs

open-source Agda library

but still WIP

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